

# How Strong is the Weak Axiom?

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# Introduction

- ▶ For sufficiently rich choice/demand data sets, possible inconsistencies with preference maximization are characterized by simple pairwise, or local consistency conditions.
- ▶ This paper is concerned with what constitutes a 'rich enough' collection of observations for such results, *independent of the specific choices observed*.

# Motivation

- ▶ *Experimental*: What experiments preclude the existence of testable implications of preference transitivity beyond pairwise-coherent choices?
- ▶ *Applied Theory*:
  - ▶ *Behavioral*: How far can full-domain assumptions in choice models be relaxed?
  - ▶ *Measurement of Irrationality*: How does structure of choice problem create potential dependencies between choice cycles?
- ▶ *Computational*: When does it suffice to verify the absence of choice cycles by checking only for 2-cycles?

# Abstract Choice I

A **choice problem** is a pair  $(X, \Sigma)$ , where:

- ▶  $X$  is a set of **alternatives**.
- ▶  $\Sigma \subseteq 2^X \setminus \{\emptyset\}$  is a collection of non-empty subsets of  $X$  called **budgets**.

These budgets correspond to the subsets of  $X$  from which we observe the agent choose.

- ▶ Assumptions on  $\Sigma$  are assumptions on *observability*.

## Abstract Choice II

A **choice correspondence** is a map  $c : \Sigma \rightarrow 2^X \setminus \{\emptyset\}$  satisfying:

$$(\forall B \in \Sigma) \quad c(B) \subseteq B.$$

The **revealed preference pair** associated to  $c$ , denoted  $(\succsim_c, \succ_c)$ , is defined via:

- ▶  $x \succsim_c y$  if there exists a budget  $B \in \Sigma$  such that  $\{x, y\} \in B$ , and  $x \in c(B)$ .
- ▶  $x \succ_c y$  if there exists a budget  $B \in \Sigma$  such that  $\{x, y\} \in B$ ,  $x \in c(B)$ , and  $y \notin c(B)$ .

## Rationalizable Choice

A choice correspondence  $c$  is **strongly rationalizable** if there exists a weak order  $\succeq$  on  $X$  such that:

$$(\forall B \in \Sigma) \quad c(B) = \{x \in B : \forall y \in B, x \succeq y\}$$

## The -ARPs

- ▶ A choice correspondence satisfies the **weak** axiom of revealed preference (WARP) if it makes no choice *reversals*:

$$x \succsim_c y \implies y \not\prec_c x.$$

- ▶ It satisfies the **generalized** axiom of revealed preference (GARP) if it contains no finite choice *cycles* of the form:

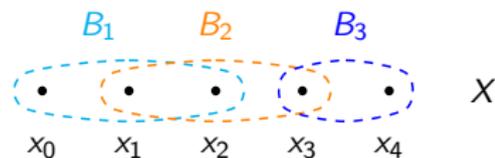
$$x_0 \succsim_c x_1 \succsim_c \cdots \succsim_c x_{N-1} \succ_c x_0.$$

# The Budget Graph

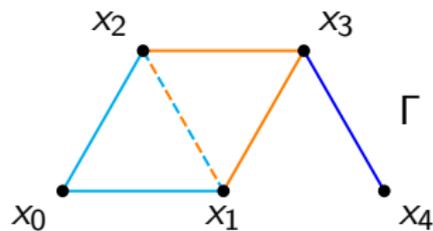
For a choice problem  $(X, \Sigma)$  its **budget graph**  $\Gamma$  is the undirected graph with vertex set  $V_\Gamma = X$  and edge set:

$$E_\Gamma = \left\{ \{x, y\} \subseteq X : \exists B \in \Sigma \text{ s.t. } \{x, y\} \subseteq B \right\}.$$

# An Example



(a) A choice problem.



(b) The budget graph.

**Figure:** A choice problem with five alternatives and three budgets.

## Cyclic Collections

For a loop  $\gamma = (V_\gamma, E_\gamma)$ , a collection of budgets  $\mathcal{B}_\gamma \subseteq \Sigma$  is a **cyclic collection** for  $\gamma$  if:

$$(\forall e \in E_\gamma) (\exists B \in \mathcal{B}_\gamma) \quad e \subseteq B.$$

A cyclic collection  $\mathcal{B}_\gamma$  for a loop  $\gamma$  is **covered** if there exists a budget  $\tilde{B} \subseteq \cup_{\tilde{B} \in \mathcal{B}_\gamma} \tilde{B}$  that either:

- (i) Contains  $V_\gamma$ ; or
- (ii) Contains a pair of elements of  $V_\gamma$  that are not connected by an edge in  $E_\gamma$ .

*Note:* Condition (i) implies (ii) if and only if  $|V_\gamma| > 3$ .

## Well-covered Budget Collections

A budget collection  $\Sigma$  is **well-covered** if, for every loop  $\gamma$  in its budget graph, every cyclic collection for  $\gamma$  is covered.

### Theorem

*Let  $(X, \Sigma)$  be a choice problem. The weak axiom of revealed preference is characteristic of strong rationalizability if and only if  $\Sigma$  is well-covered.*

# Discussion

Well-coveredness characterizes when choice cycles imply choice reversals.

It is a *richness* condition, rather than a measure of data set size:

- ▶ Adding budgets to a well-covered budget collection may create uncovered cyclic collections.
- ▶ Removing budgets may remove uncovered cyclic collections.

## Intro to Integrability

The problem of integrability of a demand function is parallel to that of the strong rationalizability of a choice correspondence.

Hurwicz, Uzawa et al. tell us a (nice) demand arises from constrained-optimal choice according to a (nice) utility if and only if its Slutsky matrix is:

- ▶ Negative semi-definite  $\iff$  weak axiom; and
- ▶ Symmetric  $\iff$  locally integrable.

*Complete domain:* we know  $x$  for every  $(p, w)$ .

# Domains

Let  $(X, \Sigma)$  be a choice problem with budget graph  $\Gamma = (X, E_\Gamma)$ .

- ▶ The **domain** associated to the choice problem is the triple  $(X, E_\Gamma, T_\Gamma)$  where:

$$T_\Gamma = \{\{x, y, z\} \subseteq X : \{x, y\}, \{y, z\}, \{x, z\} \in E_\Gamma\}.$$

- ▶ Given some collection  $\tilde{T} \subseteq T_\Gamma$ , the **subdomain** generated by  $\tilde{T}$  is the collection of 1-, 2-, and 3-element subsets of elements of  $\tilde{T}$ .

## Local Rationalizability

A choice correspondence is **locally rationalizable** if there exists an order extension  $(\succeq, \succ)$  of  $(\succsim_c, \succ_c)$  such that:

$$(\forall \tau \in T_\Gamma) \quad \succeq|_\tau \text{ is complete and transitive.}$$

## Simple (Sub)domains

A (sub)domain  $(\tilde{X}, \tilde{E}, \tilde{T})$  is **simple** if it is:

- (i) *Combinatorially Trivial*: If  $\tau, \tau' \in \tilde{T}$ , there is a unique, finite sequence of distinct elements of  $\tilde{T}$ :

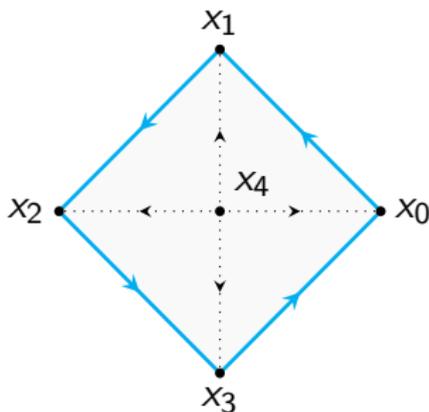
$$\tau = \tau_0, \tau_1, \dots, \tau_k = \tau'$$

such that  $\tau_j$  and  $\tau_{j+1}$  share precisely a pair of elements.

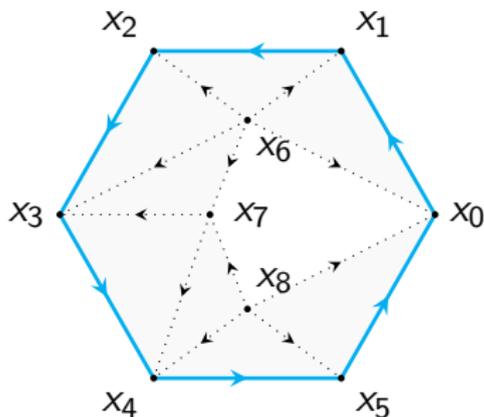
- (ii) *Topologically Trivial*: The (sub)domain has a first Betti number of zero.

By abuse of notation, we say a *domain* is simple if every loop in  $\Gamma$  is contained in a simple subdomain.

# The Triviality Conditions



**(a)** A locally rational binary relation on a topologically trivial domain that is combinatorially non-trivial.



**(b)** A locally rational binary relation on a combinatorially trivial domain that is topologically non-trivial.

**Figure:** On non-simple (sub)domains, locally rational binary relations may not be acyclic.

# Ordinal Integrability

## Theorem

*Let  $(X, \Sigma)$  be a choice problem with a simple domain. Then a choice correspondence is strongly rationalizable if and only if:*

- (i) It obeys the weak axiom; and*
- (ii) It is locally rationalizable.*

*Moreover, (i) and (ii) are equivalent to strong rationalizability if and only if the domain of  $(X, \Sigma)$  is simple.*

## Discussion

In comparison with the classical Hurwicz-Uzawa theory of integrability:

- ▶ Still require the weak axiom and a local 'no cycles' condition.
- ▶ Complete domain assumption not needed, only simplicity.
- ▶ No analytic, point-set, or order theoretic assumption on primitives required.

# Integrability Within the Empirical Content of Choice

*Tradeoff:* in classical theory, extra assumptions yield nice structure on rationalizing preference. Here only a weak order.

- ▶ The price for a theory that imposes no assumptions that are non-falsifiable by finite data sets.

## A Decomposition of Well-coveredness

Integrability theory provides an *economic* interpretation of well-coveredness.

- ▶ *Local*: Any revealed preference pair obeying the weak axiom is locally rationalizable.
  - ▶ “Negative semi-definite Slutsky matrix  $\implies$  symmetric.”
- ▶ *Global*: The domain of the choice problem is simple, hence these local conditions characterize strong rationalizability.
  - ▶ The set of observations is sampled “richly” enough.

# Conclusions

- ▶ The objective of this paper is to weaken the informational requirements of choice theory.
- ▶ A great deal of economic theory, both classical and modern, supposes an exhaustive collection of observations.
- ▶ This is a first step in studying just how widely applicable 'full domain' theories are in practice, and how they may be adapted to hold as widely as possible.

Thank you!

Any Questions?